

## LECTURE NOTES SIMULATION

## CHAPTER 8

### Random Variate Generation (cont.)

Having obtained random numbers that are Uniform(0,1) and independent, we'll learn transforming them to other Random Variables that have specific distributions.

#### (I) INVERSE TRANSFORM TECHNIQUE

Step 1. Compute the cdf of the desired R.V.  $X$ :  $F(X)$

Step 2. Set  $F(X) = R$  on the range of  $X$

Step 3. Solve the equation  $F(X) = R$  in terms of  $X$ :  $X = F^{-1}(R)$

Step 4. Given  $R_1, R_2, \dots, R_N$  : obtain  $X_i = F^{-1}(R_i)$

### CONTINUOUS DISTRIBUTIONS:

#### Exponential Distribution

$X \sim \text{Exponential}(\lambda)$

$$f(x) = \begin{cases} \lambda \cdot e^{-\lambda x} & x \geq 0 \\ 0 & \text{o.w.} \end{cases} \quad E(X) = \frac{1}{\lambda}$$

$$\text{Var}(X) = \frac{1}{\lambda^2}$$

$$F(x) = \int_{-\infty}^x f(w) dw = \int_0^x f(w) dw = \int_0^x (\lambda \cdot e^{-\lambda w}) dw$$

$$= \lambda \cdot \left[ -\frac{1}{\lambda} \cdot e^{-\lambda w} \right]_0^x = -(e^{-\lambda x} - 1) = 1 - e^{-\lambda x} \quad x \geq 0$$

$$F(x) = \begin{cases} 1 - e^{-\lambda x} & x \geq 0 \\ 0 & x < 0 \end{cases}$$

$$F(x) = \beta$$

$$1 - e^{-\lambda x} = R$$

$$1 - R = e^{-\lambda x}$$

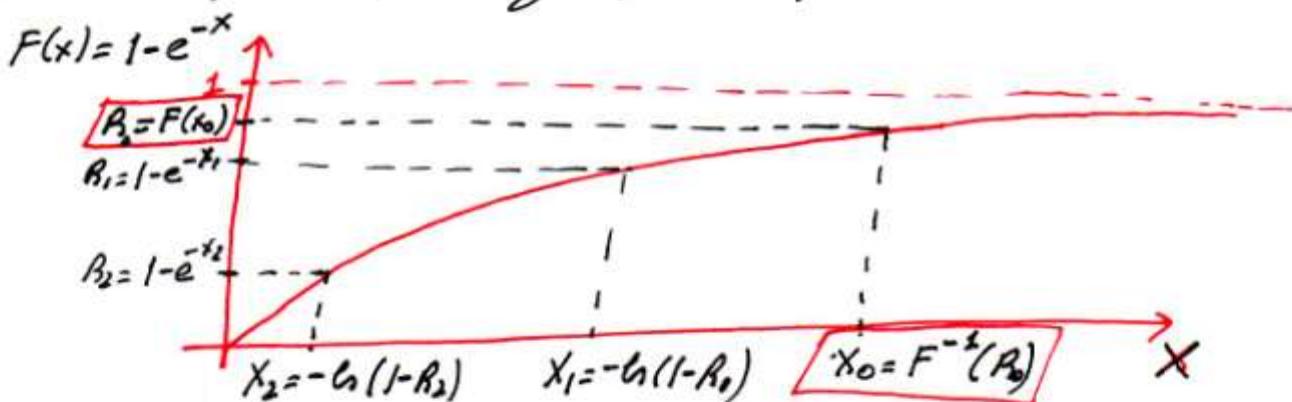
$$\ln(1-R) = -\lambda x$$

$$x = F^{-1}(R) = -\frac{1}{\lambda} \cdot \ln(1-R)$$

Note that,  $1-R \sim \text{Uniform}(0;1)$  because  $R \sim \text{Uniform}(0;1)$   
 So, we can say in short;

$$x = F^{-1}(R) = -\frac{1}{\lambda} \cdot \ln R$$

Examine the logic of Inverse transform technique  
 by the help of following graph; for  $\lambda=1$



17. Lead times have been found to be exponentially distributed with mean 3.7 days. Generate five random lead times from this distribution.

$$17) \text{ Mean} = 3,7 = \frac{1}{\lambda} \Rightarrow x_i = -3,7 \cdot \ln R_i$$

let the Random numbers be given (or from the table):

i	1	2	3	4	5
$R_i$	0,2031	0,8211	0,0390	0,6207	0,8520
$x_i$	$-3,7 \cdot \ln 0,2031$ = 5,898	$-3,7 \cdot \ln 0,8211$ = 0,725	$-3,7 \cdot \ln 0,0390$ = 12,004	$-3,7 \cdot \ln 0,6207$ = 1,765	$-3,7 \cdot \ln 0,8520$ = 0,593

1. Develop a random-variate generator for a random variable  $X$  with the pdf

$$f(x) = \begin{cases} e^{2x}, & -\infty < x \leq 0 \\ e^{-2x}, & 0 < x < \infty \end{cases}$$

4)  $F(x) = \int_{-\infty}^x f(\omega) d\omega$

$$\text{for } -\infty < x \leq 0 : F(x) = \int_{-\infty}^x e^{2\omega} d\omega = \frac{e^{2\omega}}{2} \Big|_{-\infty}^x = \frac{1}{2} (e^{2x} - e^{-\infty}) = \frac{e^{2x}}{2}$$

$$F(0) = \frac{e^0}{2} = \frac{1}{2}$$

$$\text{for } 0 < x < \infty : F(x) = \frac{1}{2} + \int_0^x e^{-2\omega} d\omega = \frac{1}{2} - \frac{1}{2} \cdot e^{-2\omega} \Big|_0^x = \frac{1}{2} - \frac{1}{2} [e^{-2x} - 1] \\ = \frac{1}{2} + \frac{1}{2} - \frac{1}{2} e^{-2x} = 1 - \frac{1}{2} \cdot e^{-2x}$$

$$F(x) = \begin{cases} \frac{1}{2} e^{2x} & -\infty < x \leq 0 \\ 1 - \frac{1}{2} e^{-2x} & 0 < x < \infty \end{cases}$$

Since 0 is a break-point and  $F(0) = \frac{1}{2}$ ; we have:

$$0 \leq R \leq \frac{1}{2} : F(x) = R \quad \frac{1}{2} < R \leq 1 : F(x) = R \\ \frac{1}{2} e^{2x} = R \quad 1 - \frac{1}{2} e^{-2x} = R$$

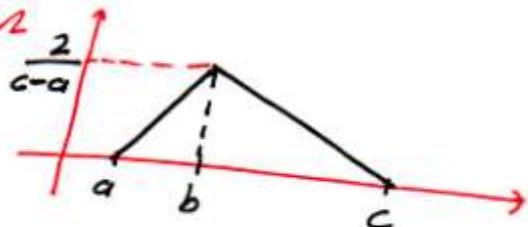
$$e^{2x} = 2R \quad 2 \cdot (1-R) = e^{-2x} \\ x = \frac{1}{2} \ln(2R) \quad x = -\frac{1}{2} \ln(2 \cdot (1-R))$$

$$x = \begin{cases} \frac{1}{2} \ln(2R) & 0 \leq R \leq 0,5 \\ -\frac{1}{2} \ln[2(1-R)] & 0,5 < R < 1 \end{cases}$$

## Triangular Distribution

$X \sim \text{Triangular}(a; b; c)$

$$f(x) = \begin{cases} \frac{2(x-a)}{(b-a)(c-a)} & a \leq x < b \\ \frac{2(c-x)}{(c-b)(c-a)} & b \leq x \leq c \\ 0 & \text{o.w.} \end{cases}$$



$$E(X) = \frac{a+b+c}{3}$$

$$\text{Mode} = b = 3E(X) - (a+c)$$

Note that triangular distribution pdf is usual line equation. Height  $\frac{2}{c-a}$  make total Area = 1.

Bdf:  $F(x)$  is given by following formula but we can also calculate it as in last exercise in a usual manner.

$$F(x) = \begin{cases} 0 & x \leq a \\ \frac{(x-a)^2}{(b-a)(c-a)} & a < x \leq b \\ 1 - \frac{(c-x)^2}{(c-b)(c-a)} & b < x \leq c \\ 1 & x > c \end{cases}$$

2. Develop a generation scheme for the triangular distribution with pdf

$$f(x) = \begin{cases} \frac{1}{2}(x-2), & 2 \leq x \leq 3 \\ \frac{1}{2}\left(2 - \frac{x}{3}\right), & 3 < x \leq 6 \\ 0, & \text{otherwise} \end{cases}$$

Generate 10 values of the random variate, compute the sample mean, and compare it to the true mean of the distribution.

3. Develop a generator for a triangular distribution with range (1, 10) and mode at  $x = 4$ .  
 4. Develop a generator for a triangular distribution with range (1, 10) and a mean of 4.

2) I'll NOT use formula of  $F(x)$ , <sup>but</sup> drive  $F(x)$  manually :)

$$2 \leq x \leq 3: F(x) = \int_{2}^x \frac{1}{2}(\omega - 2) d\omega = \frac{1}{2} \left[ \frac{\omega^2}{2} - 2\omega \right]_2^x = \frac{1}{2} \left[ \frac{x^2}{2} - 2x - \left( \frac{4}{2} - 4 \right) \right]$$

$$= \frac{1}{2} \left( \frac{x^2 - 4x + 4}{2} \right) = \frac{(x-2)^2}{4}; F(3) = \frac{(3-2)^2}{4} = \frac{1}{4}$$

$$3 < x \leq 6: F(x) = \frac{1}{4} + \int_{3}^x \frac{1}{2} \left( 2 - \frac{\omega}{3} \right) d\omega = \frac{1}{4} + \frac{1}{2} \left( 2\omega - \frac{\omega^2}{6} \right)_3^x$$

$$= \frac{1}{4} + \frac{1}{2} \left[ 2x - \frac{x^2}{6} - \left( 6 - \frac{9}{6} \right) \right] = \frac{1}{4} + \frac{1}{2} \left[ \frac{12x - x^2 - 27}{6} \right]$$

$$= \frac{1}{4} + \frac{1}{2} \left( - \frac{(x^2 - 12x + 36) + 9}{6} \right) = \frac{1}{4} + \frac{1}{2} \left( \frac{3}{2} - \frac{(x-6)^2}{6} \right) = 1 - \frac{(x-6)^2}{12}$$

$$F(x) = \begin{cases} 0 & x < 2 \\ \frac{(x-2)^2}{4} & 2 \leq x \leq 3 ; F(3) = \frac{1}{4} \\ 1 - \frac{(x-6)^2}{12} & 3 < x \leq 6 \\ 1 & x > 6 \end{cases}$$

$$0 \leq R \leq \frac{1}{4}: F(x) = R$$

$$\frac{(x-2)^2}{4} = R$$

$$(x-2)^2 = 4R$$

$$x = 2 + \sqrt{4R}$$

$$\frac{1}{4} \leq R \leq 1: F(x) = R$$

$$1 - \frac{(x-6)^2}{12} = R$$

$\underline{12(1-R)} = \underline{(6-x)^2}$   $\Rightarrow (x-6)^2$  but  
x should be within the range

$$x = 6 - \sqrt{12(1-R)}$$

$$x = \begin{cases} 2 + \sqrt{4R} & 0 \leq R \leq 0,25 \\ 6 - \sqrt{12(1-R)} & 0,25 < R \leq 1 \end{cases}$$

3)  $a = 1; b = 4; c = 10$   $(b-a)(c-a) = 3 \cdot 9 = 27$  and  
 $(c-b)(c-a) = 6 \cdot 9 = 54$

$$F(x) = \begin{cases} 0 & x \leq 1 \\ \frac{(x-1)^2}{27} & 1 < x \leq 4 \\ 1 - \frac{(10-x)^2}{54} & 4 < x \leq 10 \\ 1 & x > 10 \end{cases}$$

$$0 \leq R \leq \frac{1}{3}: F(X) = R$$

$$\frac{(x-1)^2}{27} = R$$

$$x = 1 + \sqrt{27R}$$

$$\frac{1}{3} \leq R \leq 1: F(x) = R$$

$$1 - \frac{(10-x)^2}{54} = R$$

$$54(1-R) = (10-x)^2$$

$$x = 10 - \sqrt{54(1-R)}$$

$$X = \begin{cases} 1 + \sqrt{27R} & 0 \leq R \leq 0,3333 \\ 10 - \sqrt{54(1-R)} & 0,3333 < R \leq 1 \end{cases}$$

4) Mode =  $b = 3 \cdot 4 - (10 - 1) = 12 - 9 = 3$

Apply the same procedure in §.3 with  $a=1; b=3; c=10$

### Empirical Continuous Distributions

If we have no theoretical distribution that provides a good model for the input data for a continuous random variable, we rank the observations, plot the points of Empirical cdf  $S_N(x)$  and combine the plotted points to approximate  $F(x)$  linearly. Next are two examples:

**Example:** Five observations of fire-crew response times to incoming alarms have been collected to be used in a simulation investigating possible alternative staffing and crew-scheduling policies. The data are:

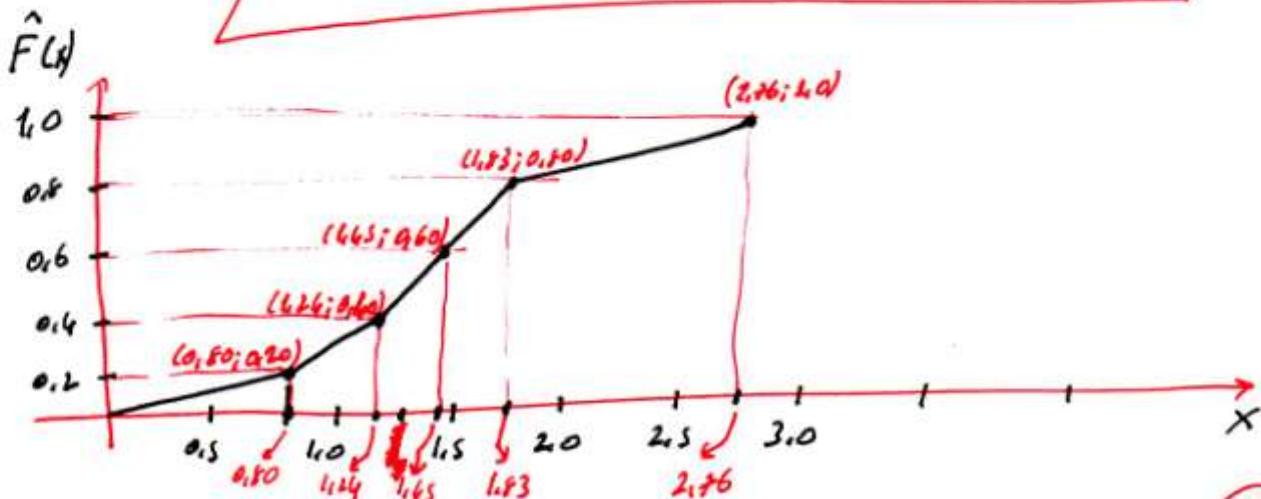
2,76    1,83    0,80    1,65    1,24

Procedure

i	Interval $X_{(i-1)} < X \leq X_{(i)}$	Probability $\hat{F}(x)$	Cum. Prob. $F(x)$	Slope $a_i = \frac{x_i - x_{(i-1)}}{\hat{F}(x_i) - \hat{F}(x_{(i-1)})}$
1	$0,00 < X \leq 0,80$	0,2	0,2	$\frac{0,80 - 0,00}{0,2} = 4,00$
2	$0,80 < X \leq 1,24$	0,2	0,4	$\frac{1,24 - 0,80}{0,2} = 2,20$
3	$1,24 < X \leq 1,65$	0,2	0,6	$\frac{1,65 - 1,24}{0,2} = 1,05$
4	$1,65 < X \leq 1,83$	0,2	0,8	$\frac{1,83 - 1,65}{0,2} = 1,90$
5	$1,83 < X \leq 2,76$	0,2	1,0	$\frac{2,76 - 1,83}{0,2} = 4,65$

Note that  $a_i$  is the slope of the  $i^{th}$  interval's straight line

$$X = \hat{F}^{-1}(R) = X_{(i-1)} + a_i \cdot \left( R - \frac{(i-1)}{n} \right)$$



So, the linear equations for  $i^{th}$  interval, namely the Empirical Random Variate generator is as follows:

$$X = \begin{cases} = 4R & 0 < R \leq 0,20 \\ = 0,80 + 2,20 \cdot (R - 0,2) & 0,20 < R \leq 0,40 \\ = 1,24 + 1,05 \cdot (R - 0,4) & 0,40 < R \leq 0,60 \\ = 1,65 + 1,90 \cdot (R - 0,6) & 0,60 < R \leq 0,80 \\ = 1,83 + 4,65 \cdot (R - 0,8) & 0,80 < R \leq 1,00 \end{cases}$$

11. Data have been collected on service times at a drive-in bank window at the Shady Lane National Bank. This data are summarized into intervals as follows:

$a_i$	Interval (Seconds)	Frequency	Rel Freq.	Cum. Freq.	slope $a_i$
$15 \div 0,07$	15-30	10	0,07	0,07	214,3
$15 \div 0,13$	30-45	20	0,13	0,20	115,4
$15 \div 0,17$	45-60	25	0,17	0,37	88,2
$30 \div 0,23$	60-90	35	0,23	0,60	130,4
$30 \div 0,20$	90-120	30	0,20	0,80	150,0
$60 \div 0,13$	120-180	20	0,13	0,93	461,5
$120 \div 0,07$	180-300	10	0,07	1,00	1714,3
		150			

Set up a table like Table 8.3 for generating service times by the table-lookup method, and generate five values of service time, using four-digit random numbers.

$$X = \begin{cases} 15 + 214,3 \cdot R & 0 < R \leq 0,07 \\ 30 + 115,4 (R - 0,07) & 0,07 < R \leq 0,20 \\ 45 + 88,2 \cdot (R - 0,20) & 0,20 < R \leq 0,37 \\ 60 + 130,4 \cdot (R - 0,37) & 0,37 < R \leq 0,60 \\ 90 + 150,0 \cdot (R - 0,60) & 0,60 < R \leq 0,80 \\ 120 + 461,5 \cdot (R - 0,80) & 0,80 < R \leq 0,93 \\ 180 + 1714,3 \cdot (R - 0,93) & 0,93 < R \leq 1,00 \end{cases}$$

10. Times to failure for an automated production process have been found to be randomly distributed with a Weibull distribution with parameters  $\beta = 2$  and  $\alpha = 10$ . Derive Equation (8.6), and then use it to generate five values from this Weibull distribution, using five random numbers taken from Table A.1.

10)

## Weibull Distribution

$$X \sim \text{Weibull}(\alpha; \beta)$$

$$f(x) = \begin{cases} \frac{\beta}{\alpha^\beta} x^{\beta-1} e^{-(x/\alpha)^\beta} & x \geq 0 \\ 0 & \text{o.w.} \end{cases}$$

$$F(x) = \int_0^x f(w) dw = \int_0^x \frac{\beta}{\alpha^\beta} \cdot w^{\beta-1} e^{-(w/\alpha)^\beta} dw$$

$$(w/\alpha)^\beta = u$$

$$F(x) = \int_0^x e^{-u} du = -e^{-u} \Big|_0^{(x/\alpha)^\beta} = 1 - e^{-(x/\alpha)^\beta} \quad x \geq 0$$

$$F(x) = R$$

$$1 - e^{-(x/\alpha)^\beta} = R$$

$$1 - R = e^{-(x/\alpha)^\beta}$$

$$\ln(1-R) = -(x/\alpha)^\beta$$

$$[-\ln(1-R)]^{1/\beta} = +\frac{x}{\alpha}$$

$$x = +\alpha \cdot [-\ln R]^{1/\beta}$$

Since R and 1-R has same distribution

$$\beta=2; \alpha=10 \Rightarrow x_i = +10 \cdot [-\ln R_i]^{1/2}$$

i	1	2	3	4	5
R <sub>i</sub>	0,8252	0,7926	0,11257	0,1093	0,6756
X <sub>i</sub>	4,383	4,821	14,601	14,878	8,621

$$\hookrightarrow 10[-\ln 0,8252]^{1/2} - - -$$



19. A machine is taken out of production either if it fails or after 5 hours, whichever comes first. By running similar machines until failure, it has been found that time to failure,  $X$ , has the Weibull distribution with  $\alpha = 8$ ,  $\beta = 0.75$ , and  $v = 0$  (refer to Sections 5.4 and 8.1.3). Thus, the time until the machine is taken out of production can be represented as  $Y = \min(X, 5)$ . Develop a step-by-step procedure for generating  $Y$ .

19)   
 Step 0  $i=1$ , Set  $N$   
 Step 1 Generate  $R_i$   
 Step 2  $X_i = 8 \cdot [-\ln R_i]^{0.75}$   
 Step 3 If  $X_i > 5$  then  $Y_i = X_i$   
 o.w.  $Y_i = 5$   
 Step 4 If  $N=i$ , stop.  
 o.w.  $i \leftarrow i+1$  and GOTO step 1

### Standard Normal Distribution

#### Schmeiser Method

$$Z = \frac{R^{0.135} - (1-R)^{0.135}}{0.1975}$$

#### Box and Muller Method

$$Z_1 = (-2 \ln R_1)^{1/2} \cos(2\pi R_2)$$

$$Z_2 = (-2 \ln R_1)^{1/2} \sin(2\pi R_2)$$

Note That  $X_i = \mu + z_i \sigma \sim \text{Normal}(\mu; \sigma^2)$

18. Regular maintenance of a production routine has been found to vary and has been modeled as a normally distributed random variable with mean 33 minutes and variance 4 minutes<sup>2</sup>. Generate five random maintenance times from the given distribution.

18) Schmeiser $i$	1	2	3	4	5
$R_i$	0,3096	0,7776	0,0292	0,0152	0,2634
$Z_i$	-0,494	0,761	-0,901	-2,176	-0,630
$33 + 2Z_i = X_i$	32,01	34,52	29,20	28,65	31,74

## Continuous Uniform Distribution

\* Remember,  $X \sim \text{Uniform}(a; b)$

$$f(x) = \begin{cases} \frac{1}{b-a} & a \leq x \leq b \\ 0 & \text{o.w.} \end{cases} \quad F(x) = \begin{cases} 0 & x < a \\ \frac{x-a}{b-a} & a \leq x \leq b \\ 1 & x > b \end{cases}$$

Then;  $F(x) = R$

$$\frac{x-a}{b-a} = R$$

$\boxed{x = a + (b-a) \cdot R}$  generates Uniform Random Numbers over  $a$  and  $b$  from  $R$

Example Generate 5 Random Numbers  $X \sim \text{Uniform}[-11, 17]$

Answer

$$F(x) = \begin{cases} 0 & x < -11 \\ \frac{x+11}{28} & -11 \leq x \leq 17 \\ 1 & x > 17 \end{cases}$$

$$F(x) = R$$

$$\frac{x+11}{28} = R \Rightarrow \underline{\underline{x = 28R - 11}}$$

i	1	2	3	4	5
$R_i$	0,3045	0,0674	0,4683	0,5122	0,2620
$X_i$	-2,47	-9,67	1,55	3,34	-3,66
	$= 28 \cdot 0,3045 - 11$	-	-	-	

## DISCRETE DISTRIBUTIONS:

Empirical Discrete Distribution;

**Example 4** At the end of any day, the # of shipments on a loading dock company is either 0, 1 or 2 with observed relative frequency of occurrence of 0,50; 0,30; 0,20 respectively. Simulate 10 shipments and estimate the mean shipment/day. Compare it with the real mean.

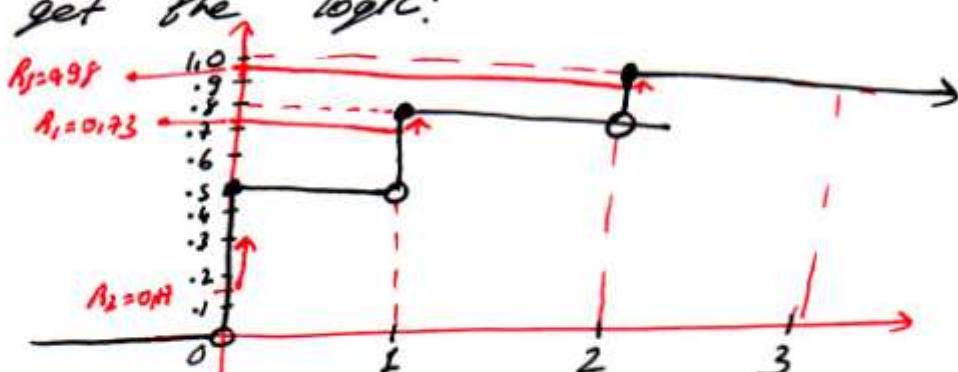
**Answer**

x	p(x)	F(x)
0	0,50	0,50
1	0,30	0,80
2	0,20	1,00

$$F(x) = \begin{cases} 0 & x < 0 \\ 0,5 & 0 \leq x < 1 \\ 0,8 & 1 \leq x < 2 \\ 1,0 & x \geq 2 \end{cases}$$

Then,  $F(x) = \beta \Rightarrow X = F^{-1}(\beta)$  gives;

Observe the following graph to get the logic:



$$\hat{\mu} = \bar{X} = \frac{\sum X_i}{n} = \frac{7}{10} = 0,70$$

$$\mu = E(X) = 0,50 \cdot 0 + 0,30 \cdot 1 + 0,20 \cdot 2 = 0,50$$

$$X = \begin{cases} 0 & \beta \leq 0,5 \\ 1 & 0,5 < \beta \leq 0,8 \\ 2 & 0,8 < \beta \leq 1,0 \end{cases}$$

i	R <sub>i</sub>	X <sub>i</sub>
1	0,73	1
2	0,17	0
3	0,98	2
4	0,30	0
5	0,04	0
6	0,44	0
7	0,51	1
8	0,26	0
9	0,75	1
10	0,97	2

$\sum X_i = 7$

## Discrete Uniform Distribution

remember;  $X \sim \text{Discrete Uniform}(k)$

$$p(x) = \frac{1}{k}, \quad x = 1, 2, \dots, k$$

Then;  $F(X) = R$  implies

$$F(x_{i+1}) = r_{j+1} \leq \beta \leq r_i = F(x_i)$$

$$r_{i-1} = \frac{i-1}{k} < \beta \leq r_i = \frac{i}{k}$$

$$j-t < \beta k \leq j$$

$$Bk \leq i < Bk+1 \Rightarrow$$

$$F(x) = \begin{cases} 0 & x < 1 \\ \frac{1}{k} & 1 \leq x < 2 \\ \frac{2}{k} & 2 \leq x < 3 \\ \vdots & \\ \frac{k-1}{k} & k-1 \leq x < k \\ 1 & k \leq x \end{cases}$$

→ Round UP fractions!

$$X = [Rk]$$

13. For a preliminary version of a simulation model, the number of pallets,  $X$ , to be loaded onto a truck at a loading dock was assumed to be uniformly distributed between 8 and 24. Devise a method for generating  $X$ , assuming that the loads on successive trucks are independent. Use the technique of Example 8.5 for discrete uniform distributions. Finally, generate loads for 10 successive trucks by using four-digit random numbers.

13) Note that,  $y \sim \text{Discrete Uniform } (17)$   $\rightarrow 26 - 8 + 1 = 17 \text{ Number}$

Then,  $X = Y + 7 \rightarrow 8-1=7$  because  $Y = 1, 2, \dots, 17$   
 $X = 8, 9, \dots, 24$

*Step D Set  $i = 1$*

*Step 1* generate  $R_j$

$$\text{Step 2 } x_i = \lceil 17 \cdot R_i \rceil$$

$$\text{Step 3} \quad y_i = x_i + 7$$

*Step 4 If  $N=i$  stop.*

O.W.  $i \leftarrow i+1$  and  
GOTO Step 1

	$i$	1	2	3	4	5
$R_i$	0,8832	0,7686	0,9691	0,3026	0,8670	
$X_i$	16	$\overbrace{17, \text{ 0,88327}}^{\rightarrow}$ - - -	13	17	6	15
$y_i$	23	$\overbrace{16+7}^{\rightarrow}$ - - -	20	24	13	22
	$i$	6	7	8	9	10
$A_i$	0,0532	0,2143	0,1111	0,5566	0,6113	
$X_i$	1	4	2	10	11	
$y_i$	8	11	9	17	18	(18)

9. The cdf of a discrete random variable  $X$  is given by

$$F(x) = \frac{x(x+1)(2x+1)}{n(n+1)(2n+1)}, \quad x = 1, 2, \dots, n$$

When  $n = 4$ , generate three values of  $X$ , using  $R_1 = 0.83$ ,  $R_2 = 0.24$ , and  $R_3 = 0.57$ .

$$9) \quad n=4 \Rightarrow F(x) = \frac{x(x+1)(2x+1)}{4 \cdot 5 \cdot 9} \quad x = 1, 2, 3, 4$$

$$F(1) = \frac{1 \cdot 2 \cdot 3}{4 \cdot 5 \cdot 9} = 0.033 \quad F(2) = \frac{2 \cdot 3 \cdot 5}{4 \cdot 5 \cdot 9} = 0.167 \quad F(3) = \frac{3 \cdot 6 \cdot 7}{4 \cdot 5 \cdot 9} = 0.667 \quad F(4) = 1$$

$$F(x) = \begin{cases} 0 & x < 1 \\ 0.033 & 1 \leq x < 2 \\ 0.167 & 2 \leq x < 3 \\ 0.667 & 3 \leq x < 4 \\ 1 & x \geq 4 \end{cases} \Rightarrow X = \begin{cases} 1 & R \leq 0.033 \\ 2 & 0.033 < R \leq 0.167 \\ 3 & 0.167 < R \leq 0.667 \\ 4 & R > 0.667 \end{cases}$$

$i$	1	2	3
$R_i$	0.83	0.24	0.57
$x_i$	4	3	4

### Geometric Distribution

$$X \sim \text{Geometric } (\rho) \quad E(X) = \frac{1}{\rho}$$

$$p(x) = \rho(1-\rho)^x, \quad x = 0, 1, 2, \dots$$

(Usually  $x$  starts from 1 but we have started it from 0 for ease of calculations)

$$F(x) = \sum_{w=0}^x \rho(1-\rho)^w = \rho \cdot \sum_{w=0}^x (1-\rho)^w = \rho \cdot \underbrace{\frac{1 - (1-\rho)^{x+1}}{1 - (1-\rho)}}_{\hookrightarrow \text{Finite sum formula}} = 1 - (1-\rho)^{x+1}$$

$$F(x) = 1 - (1-\rho)^{x+1}$$

$$F(x-1) = 1 - (1-p)^x < \beta \leq 1 - (1-p)^{x+1} = F(x)$$

$$(1-p)^{x+1} \leq 1 - \beta < (1-p)^x$$

$$(x+1) \ln(1-p) \leq \ln(1-\beta) < x \cdot \ln(1-p)$$

$$x+1 \geq \frac{\ln(1-\beta)}{\ln(1-p)} \text{ and } x < \frac{\ln(1-\beta)}{\ln(1-p)}$$

$$\text{So; } \frac{\ln(1-\beta)}{\ln(1-p)} - 1 \leq x < \frac{\ln(1-\beta)}{\ln(1-p)}$$

$$x = \left\lceil \frac{\ln(1-\beta)}{\ln(1-p)} - 1 \right\rceil$$

When we want geometric random variables that start from  $a$ , we have,  $x = a, a+1, a+2, \dots$

$$\text{and } x = a + \left\lceil \frac{\ln(1-\beta)}{\ln(1-p)} - 1 \right\rceil \quad (\text{usually } a=1 \text{ as asserted})$$

(Here, for  $a \geq 1$ , we have the mean

$$E(X) = \frac{1}{p} + a - 1$$

**Example** Generate 31 Random Variables with mean 4 on the Range  $\{x \geq 1\}$

$$\frac{1}{p} = 4 \Rightarrow p = \frac{1}{4} = 0.25 \text{ and } \frac{1}{\ln(1-p)} = \frac{1}{\ln 0.75} = -3.676$$

$$x_i = 1 + \lceil -3.676 \cdot \ln(1-R_i) - 1 \rceil$$

i	1	2	3
R <sub>i</sub>	0.932	0.105	0.687
x <sub>i</sub>	10 = 1 + [-3.676 · ln(1-0.932) - 1]	1 = --	5 = --

15. The weekly demand,  $X$ , for a slow-moving item has been found to be approximated well by a geometric distribution on the range  $\{0, 1, 2, \dots\}$  with mean weekly demand of 2.5 items. Generate 7 values of  $X$ , demand per week, using random numbers from Table A.1. (Hint: For a geometric distribution on the range  $\{q, q+1, \dots\}$  with parameter  $p$ , the mean is  $1/p + q - 1$ .)

$$15) \frac{1}{p} = 2.5 \Rightarrow p = \frac{1}{2.5} = 0.4 ; \frac{1}{q(1-p)} = \frac{1}{0.6} = 1.6666666666666667$$

$$X_i = \lceil -1.6666666666666667 \ln R_i - 1 \rceil$$

$i$	1	2	3	4	5	6	7
$R_i$	0.3157	0.0782	0.0906	0.1992	0.2992	0.2960	0.9661
$X_i$	2	4	4	3	2	2	0

$$\lceil -1.6666666666666667 \ln 0.3157 - 1 \rceil = \lceil 1.257 \rceil = 2$$

## (II) ACCEPTANCE-REJECTION TECHNIQUE

To get the idea of acceptance-rejection technique, we may start with the following simple example;

Simple Example:)

Let, we want to generate  $X \sim \text{Uniform}(0, 25; 1)$

We know that  $X_i = 0.75R_i + 0.25$  generates  $X_i$ , more efficiently. (We'll see why this is more efficient)

Alternatively, we may generate  $R_i \sim \text{Uniform}(0, 1)$  and "Accept" as  $X_i$  if  $R_i \geq 0.25$  (Namely in the range)

Step 1 Generate  $R$

Step 2 If  $R \geq 0.25$ , Accept  $X=R$ , GOTO step 3. o.w. GOTO Step 1.

Step 3 If another  $X$  is needed, GOTO step 1. o.w. STOP!

Note that;

$$P(a < R < b | 0.25 \leq R \leq 1) = \frac{P(a < R < b)}{P(0.25 \leq R \leq 1)} = \frac{b-a}{\frac{3}{4}}$$

But,  $X \sim \text{Uniform}(\frac{1}{4}, 1)$

$$F(x) = \frac{x - \frac{1}{4}}{1 - \frac{1}{4}} = \frac{x - \frac{1}{4}}{\frac{3}{4}} \quad \text{and} \quad P(a < X < b) = F(b) - F(a) = \frac{b-a}{\frac{3}{4}}$$

So, this algorithm truly generates  $x_i$ !

Now, we are interested in the efficiency of this algorithm, which depends on "Number of Rejections"

Let  $Y$ : Random Number required to generate first  $X$ :

Remember,  $Y \sim \text{Geometric}(p = \frac{3}{4}) \rightarrow p = P(0.25 \leq R \leq 1)$

$$E(Y) = \frac{1}{p} = \frac{1}{\frac{3}{4}} = \frac{4}{3}$$

Let, we need 1000 random  $x_i$ . How many  $y_i$ , on the average, we need to generate 1000  $x_i$ 's?

Let :  $N$ : Random Numbers required to generate 1000  $x_i$ .

Remember,  $N \sim \text{Negative Binomial}(r=1000; p = \frac{3}{4})$

Also Note that  $N = \sum_{i=1}^{1000} Y_i$

$$E(N) = E\left[\sum_{i=1}^{1000} Y_i\right] = \sum_{i=1}^{1000} E(Y_i) = \sum_{i=1}^{1000} \frac{4}{3} = 1000 \cdot \frac{4}{3} = 1333$$

So, we "waste" 333 Random Numbers, which decreases the performance of the algorithm. The performance may be even much worse! (i.e. if  $X \sim \text{Uniform}(0.9, 1)$ )

## Poisson Distribution

$$N \sim \text{Poisson}(\alpha)$$

$$\rho(n) = P(N=n) = \frac{e^{-\alpha} \cdot \alpha^n}{n!} \quad n=0, 1, 2, \dots$$

$$E(N) = \alpha \quad \text{and} \quad \text{Var}(N) = \alpha$$

Remember from Stochastic Processes, Number of events per unit time with rate  $\alpha$  can be well modeled with Poisson distribution.

Also remember the relationship with Poisson process and Exponential Distribution. If  $N(t)$  is a Poisson process then "T: Interarrival Time" OR "Time until Next Event" follows an Exponential Distribution with the same rate:

$$N(t) \sim \text{Poisson}(\alpha t) \Leftrightarrow T \sim \text{Exponential}(\alpha)$$

Therefore, if  $A_j$  is the time of  $j^{\text{th}}$  arrival,

$$N=n \Leftrightarrow A_1 + A_2 + \dots + A_n \leq t \leq A_1 + A_2 + \dots + A_n + A_{n+1}$$

So, we have;  $\sum_{i=1}^n A_i \leq t < \sum_{i=1}^{n+1} A_i$  but  $A_i = -\frac{1}{\alpha} \ln R_i$   
 (Generation of Exponential A.V.)

$$\sum_{i=1}^n -\frac{1}{\alpha} \ln R_i \leq t < \sum_{i=1}^{n+1} -\frac{1}{\alpha} \ln R_i$$

$$\text{or } \prod_{i=1}^n R_i \geq e^{-\alpha} > \prod_{i=1}^{n+1} R_i$$

$$\prod_{i=1}^n R_i \geq e^{-\alpha} > \prod_{i=1}^{n+1} R_i$$

This fact is the logic of the following algorithm:

**Step 1:** Set  $n=0$ ;  $P=1$

**Step 2:** Generate  $R_{n+1}$  and  $P \leftarrow P \cdot R_{n+1}$

**Step 3:** If  $P \leq e^{-\alpha}$  then accept  $N=n$

o.w. Reject current  $n$ ,  $n \leftarrow n+1$  and GOTO Step 2

Clearly, this algorithm generates a single  $N$ . We loop it as many times as needed.

Note that, to generate a single  $N=n$ , we need  $n+1$  random numbers  $R_i$ . So, on the average  $E(N+1) = \alpha + 1$  numbers are required to generate FIRST  $N$  and we must generate  $r(\alpha + 1)$  numbers to obtain  $r N's$ , which is quite inefficient if  $\alpha$  is large.

**Example** Let  $X_i \sim \text{Poisson } (\alpha=0.2)$ . Generate 3  $X_i$ 's.

First Compute  $e^{-\alpha} = e^{-0.2} = 0.8187$ .

1. $n=0, P=1$	1. $n=0, P=1$
2. $R_1 = 0.4357, P=1 \cdot R_1 = 0.4357$	2. $R_1 = 0.8353, P=1 \cdot 0.8353 = 0.8353$
3. $P=0.4357 \leq e^{-\alpha} = 0.8187$ , accept $N=0$	3. Reject $n=0, n=1$
1. $n=0, P=1$	2. $R_2 = 0.9952$
2. $R_2 = 0.6146, P=1 \cdot R_2, 0.6146$	3. $P=R_1 \cdot R_2 = 0.8313$
3. $P=0.6146 \leq e^{-\alpha} = 0.8187$ , accept $N=0$	3. Reject $n=1, n=2$
i/ 1 2 3	2. $R_3 = 0.8004$
N/ 0 0 2	3. $P=R_1 \cdot R_2 \cdot R_3 = 0.6654$
	3. Accept $N=2$

16. In Exercise 15, suppose that the demand has been found to have a Poisson distribution with mean 2.5 items per week. Generate values of  $X$ , demand per week, using random numbers from Table A.1. Discuss the differences between the geometric and the Poisson distributions.

$$16) e^{-\alpha} = e^{-2.5} = 0.082 \quad N \sim \text{Poisson}(\alpha = 2.5)$$

$n$	Step 1						
$R_{n+1}$	-	0	0	1	0	0	1
$P_{n+1} = p_n \cdot R_{n+1}$	1	0.1093	0.0520	1	0.7787	0.1648	0.6397
$A/R$	-	R	R	-	R	R	A
$n$	Step 1						
$R_{n+1}$	-	0	0	1	2	3	4
$P_{n+1} = p_n \cdot R_{n+1}$	1	0.9667	0.6791	0.7926	0.2615	0.6651	- 0.0258
$A/R$	-	R	R	R	R	A	- A

The Generated Random Numbers are 1; 2; 4 and 0.

Note that, we generated 11  $R_i$ 's to generate 4  $N_i$ 's.

Average Number of  $R_i$ 's to generate 4  $N_i$ 's is:-

$$r. E(N+1) = r. (\alpha + 1) = 4. (2.5 + 1) = \underline{\underline{14}}$$

Normal Approximation:

$$\boxed{Z = \frac{N - \alpha}{\sqrt{\alpha}} \Rightarrow N_i = [\alpha + \sqrt{\alpha} Z_i - 0.5]} \text{ can be used}$$

to generate Poisson Random Variables for  $\alpha$  is large  
(practically  $\alpha \geq 15$ )

Remember, we have seen how to generate  $Z_i$ .

## Nonstationary Poisson Process (NSPP)

Remember, a NSPP is an arrival process with an arrival rate that varies with time. The following algorithm, which is also called "thinning" algorithm can be used to generate interarrival times from a NSPP with arrival rate  $\lambda(t)$ ,  $0 \leq t \leq T$

**Step 1.** Let  $\lambda^* = \max_{0 \leq t \leq T} \lambda(t)$  and set  $t=0, i=1$

**Step 2.** Generate  $E \sim \text{Exponential}(\lambda^*)$  and  $t \leftarrow t + E$

**Step 3.** Generate  $R \sim \text{Uniform}(0; 1)$

If  $R \leq \frac{\lambda(t)}{\lambda^*}$  then  $T_i = t$  and  $i \leftarrow i+1$

**Step 4.** GOTO Step 2

This algorithm works for any integrable arrival rate function. However, we'll only see the case that Arrival rate function is a piecewise-constant function.

Example 4	Mean time between arrivals (min)	Arrival Rate: $\lambda(t)$ (arrivals/min)
0	15	$1/15$
60	12	$1/12$
120	7	$1/7$
180	5	$1/5 = \lambda^*$
240	8	$1/8$
300	10	$1/10$
360	15	$1/15$
420	20	$1/20$
480	20	$1/20$

Generate first two arrivals.

**Step 1:**  $\lambda^* = \max_{0 \leq t \leq T} \lambda(t) = \frac{1}{5}$ ;  $t=0$ ;  $i=1$

**Step 2:**  $R = 0,2130$ ;  $E = -5 \ln(0,2130) = 13,13$ ;  $t = 0 + 13,13 = 13,13$

**Step 3:**  $R = 0,8830 \not\leq \frac{\lambda(13,13)}{\lambda^*} = \frac{1/15}{1/5} = \frac{1}{3}$  do NOT generate on interval!

**Step 4:** GOTO Step 2

**Step 2:**  $R = 0,5530$ ;  $E = -5 \ln(0,5530) = 2,96$ ;  $t = 13,13 + 2,96 = 16,09$

**Step 3:**  $R = 0,0240 \leq \frac{\lambda(16,09)}{\lambda^*} = \frac{1/15}{1/5} = \frac{1}{3}$

$$\boxed{T_1 = t = 16,09} \quad i = i+1 = 2$$

**Step 4:** GOTO Step 2

**Step 2:**  $R = 0,0001$ ;  $E = -5 \ln(0,0001) = 46,05$ ;  $t = 16,09 + 46,05 = 62,14$

**Step 3:**  $R = 0,1463 \leq \frac{\lambda(62,14)}{\lambda^*} = \frac{1/12}{1/5} = \frac{5}{12} = 0,4167$

$$\boxed{T_2 = t = 62,04} \quad i = i+1 = 3$$

29. Let time  $t = 0$  correspond to 6 A.M., and suppose that the arrival rate (in arrivals per hour) of customers to a breakfast restaurant that is open 6–9 A.M. is

$$\lambda(t) = \begin{cases} 30, & 0 \leq t < 1 \\ 45, & 1 \leq t < 2 \\ 20, & 2 \leq t \leq 4 \end{cases}$$

Derive a thinning algorithm for generating arrivals from this NSPP.

29) **Step 1:**  $\lambda^* = 45$  because given numbers are directly rate.

$$E_i = -\frac{1}{45} \ln R_i \quad \begin{matrix} t=0 \\ i=1 \end{matrix} \quad \underline{0,5766} \quad \underline{0,2379} \quad \underline{0,1767} \quad \dots$$

Let  $R$  be given:  $0,7788 \quad 0,8894 \quad 0,6394 \quad 0,4755 \quad 0,0260$

For ease of calculation, we may take unit time as minutes. Then,

$$\lambda(t) = \begin{cases} \frac{30}{60} = 0.5 \text{ arr/min} & 0 \leq t < 60 \\ \frac{45}{60} = 0.75 \text{ arr/min} & 60 \leq t < 120 \\ \frac{20}{60} = 0.33 \text{ arr/min} & 120 \leq t \leq 140 \end{cases}$$

$$\lambda^* = 0.75 \quad z=0 \quad i=1 \quad E_i = -\frac{4}{3} \ln R_i$$

$$\text{Step 2: } E = -\frac{4}{3} \ln 0.7788 = 0.34 \quad t = 0.34 + 0 = 0.34$$

$$\text{Step 3: } R = 0.8894 \quad \frac{\lambda(0.34)}{\lambda^*} = \frac{0.50}{0.75} = 0.6667$$

No NOT generate an interval

Step 4 GOTO step 2

$$\text{Step 2: } E = -\frac{4}{3} \ln 0.6394 = 0.59 \quad t = 0.34 + 0.59 = 0.93$$

$$\text{Step 3: } R = 0.6755 \leq \frac{\lambda(0.93)}{\lambda^*} = 0.6667; \quad \boxed{T_1 = t = 0.93} \quad i=2$$

Step 4 GOTO step 2

$$\text{Step 2: } E = -\frac{4}{3} \ln 0.0260 = 4.85; \quad t = 0.93 + 4.85 = 5.78$$

$$\text{Step 3: } R = 0.1767 \leq \frac{\lambda(5.78)}{\lambda^*} = 0.6667; \quad \boxed{T_2 = 5.78}$$

Note that, when the period is that of  $\lambda^*$ , all arrivals will be Accepted. The idea is basically to lower the arrivals proportional to  $\lambda^*$ . Namely, an arrival will be Accepted with probability  $\frac{\lambda(t)}{\lambda^*}$ .

## (III) CONVOLUTION METHOD

The probability distribution of a sum of two or more independent random variables is called a convolution of the distribution of the original variables.

The convolution method thus refers to adding together two or more random variables to obtain a new random variable with the desired distribution.

### Erlang Distribution

$$X \sim \text{Erlang}(k, \theta)$$

$$E(X) = \frac{k}{\theta} \quad \text{Var}(X) = \frac{k}{\theta^2}$$

Remember, Erlang distribution is sum of  $k$  i.i.d Exponential Random Variables with mean  $\frac{1}{k\theta}$ .

$$\text{let } Y_i \sim \text{Exponential}(k\theta)$$

$$\text{Then; } X = \sum_{i=1}^k Y_i \quad \text{but } Y_i = -\frac{1}{k\theta} \ln R_i$$

$$X = \sum_{i=1}^k -\frac{1}{k\theta} \ln R_i = -\frac{1}{k\theta} \ln \left( \prod_{i=1}^k R_i \right)$$

### Negative Binomial Distribution

$$X \sim \text{Negative Binomial}(r; p)$$

$$E(X) = \frac{r}{p} \quad \text{Var}(X) = \frac{r(1-p)}{p^2}$$

Remember,  $Y \sim \text{Geometric}(p) \Rightarrow Y$  is the number of trials to obtain FIRST success and  $X$  is Number of trials to obtain  $r^{\text{th}}$  success.

Then, 
$$X = \sum_{i=1}^r Y_i \quad Y_i \sim \text{Geometric}(p)$$
  

$$X \sim \text{Negative Binomial}(r; p)$$

14. Develop a method for generating values from a negative binomial distribution with parameters  $p$  and  $k$ , as described in Section 5.3. Generate 3 values when  $p = 0.8$  and  $k = 2$ . [Hint: Think about the definition of the negative binomial as the number of Bernoulli trials until the  $k^{\text{th}}$  success.]

14) We have learned how to generate Geometric R.V's by Inverse transformation technique. Alternatively, we may generate them by Acceptance-Rejection Technique.

Step 0: set  $X=0$ , set  $r$ , set  $i=1$ , set  $p$

Step 1: set  $n=1$

Step 2: Generate  $R$

Step 3: If  $R \leq p$  Accept  $Y_i = n$   
 o.w., repeat, GOTO Step 2

Step 4:  $X \leftarrow X + Y_i$

Step 5: If  $i=k$ , stop

o.w.,  $i \leftarrow i+1$  and GOTO step 1

First value

$X=0$ ;  $r=2$ ;  $i=1$ ;  $p=0.8$

$n=1$

$R=0.8253 \not\leq 0.8$

$n=2$   
 $R=0.0733 \not\leq 0.8$ ;  $y=2$

$$X=0+2=2$$

$i=1 \neq 2$ , GOTO Step 1,  $i=2$

$n=1$

$R=0.7878 \leq 0.8$ ;  $y=1$

$$X=2+1=3$$

$i=2=2$  STOP